Кетје and Кеуак

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Outline









Overview

Inspired by Keccak and DUPLEX

KEYAK targeting high performances

- Using reduced-round Keccak-f[1600] or Keccak-f[800]
- Optionally parallelizable

■ KETJE targeting lightweight

■ Using reduced-round Keccak-f[400] or Keccak-f[200]

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Two approaches

Keyak:

- DUPLEXWRAP
- A (strong) permutation
 - fixed #rounds
- Block-oriented
- Cryptanalysis
 - permutation-level

KETJE:

- MONKEYWRAP
- A (thin) round function#rounds in phases
- Stream-oriented
- Cryptanalysis
 - round function + construction

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Nonce-based AE function

- 128-bit security (incl. multi-target)
- Sequence of header-body pairs
 - keeping the state during the session
- Optionally parallelizable
- Using reduced-round Keccak-*f*[1600] or Keccak-*f*[800], to allow
 - implementation re-use
 - cryptanalysis re-use
 - reasonable side-channel protections

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Κεγακ

Duplex layer



КЕССАК-p[1600, $n_r = 12$] ог КЕССАК-p[800, $n_r = 12$]

DUPLEXWRAP

- is a nonce-based authenticated encryption mode;
- works on sequences of header-body pairs.



A⁽¹⁾ contains the key and must be unique, e.g.,
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In general: $A^{(1)} = \text{key}||\text{nonce}||$ associated data.

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Keya

Inside DUPLEXWRAP



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KEYAK instances and efficiency

Name	Width <i>b</i>	Parallelism P
Ocean Keyak	1600	4
Sea Keyak	1600	2
Lake Κεγακ	1600	1
River Keyak	800	1

- Processing for LAKE KEYAK
 - Iong messages: about 50 % of SHAKE128
 - short messages: 24 rounds
- Working memory footprint
 - reasonable on high- and middle-end platforms
 - not ideal on constrained platforms

Security of KEYAK

Generic security of KEYAK thanks to a combination of results:

- Sound tree hashing modes [IJIS 2013] for parallelized modes
- Keyed sponge indistinguishability [SKEW 2011 + work in progress]
- SPONGEWRAP generic security [SAC 2011], adapted to DUPLEXWRAP

Safety margin against shortcut attacks:

- Practical attacks up to 6 rounds [Dinur et al. SHA-3 2014]
- Academic attacks up to 9 rounds [Dinur et al. SHA-3 2014]

Кетј

Outline







4 Conclusions and Current Developments

Кетје

KETJE goals

- Nonce-based AE function
- 96-bit or 128-bit security (incl. multi-target)
- Sequence of header-body pairs
 - keeping the state during the session
- Small footprint
- Target niche: secure channel protocol on secure chips
 - banking card, ID, (U)SIM, secure element, FIDO, etc.
 - secure chip has strictly incrementing counter
- Using reduced-round Keccak-*f*[400] or Keccak-*f*[200], to allow
 - implementation re-use
 - cryptanalysis re-use
 - reasonable side-channel protections

Кетје

Inside KETJE: the MONKEYDUPLEX layer



 $n_{\text{start}} = 12$ rounds should provide strong instance separation

- $n_{\text{step}} = 1, r = 2b/25$ should avoid single-instance state retrieval
- $n_{\text{stride}} = 6$ rounds should avoid a forgery with one instance

Кетј

Inside MONKEYWRAP



KETJE instances and lightweight features

feature		Ketje Jr	Ketje Sr	
state size		25 bytes	50 bytes	
block size		2 bytes	4 bytes	
processing		computational cost		
initialization	per session	12 rounds	12 rounds	
wrapping	per block	1 round	1 round	
8-byte tag comp.	per message	9 rounds	7 rounds	

Outline





3 Ketje



Current developments

Optimized software implementations

- Gross estimations can be derived from KECCAK
- LAKE KEYAK expected twice faster than SHAKE128
- There might be interesting improvement with new AVX512 (VPTERNLOG, rotations and 32 registers)
- Hardware implementations

Conclusions

Thanks for your attention!

